CS420: Operating Systems

Threading Issues

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Threading Issues

• There are a variety of issues to consider with multithreaded programming
  - Semantics of `fork()` and `exec()` system calls
  - Thread cancellation
    • Asynchronous or deferred
  - Signal handling
    • Synchronous and asynchronous
  - Thread pooling
  - Thread-specific data
    • Create facility needed for data private to thread
Semantics of \texttt{fork()} and \texttt{exec()}

- Recall that when \texttt{fork()} is called, a separate, duplicate process is created.

- How should \texttt{fork()} behave in a multithreaded program?
  - Should all threads be duplicated?
  - Should only the thread that made the call to \texttt{fork()} be duplicated?

- In some systems, different versions of \texttt{fork()} exist depending on the desired behavior.
  - Some UNIX systems have \texttt{fork1()} and \texttt{forkall()}
    - \texttt{fork1()} only duplicates the calling thread
    - \texttt{forkall()} duplicates all of the threads in a process
  - In a POSIX-compliant system, \texttt{fork()} behaves the same as \texttt{fork1()}

Semantics of \texttt{fork()} and \texttt{exec()} 

• The \texttt{exec()} system call continues to behave as expected
  - Replaces the entire process that called it, including all threads

• If planning to call \texttt{exec()} after \texttt{fork()}, then there is no need to duplicate all of the threads in the calling process
  - All threads in the child process will be terminated when \texttt{exec()} is called
  - Use \texttt{fork1()}, rather than \texttt{forkall()} if using in conjunction with \texttt{exec()}
Thread Cancellation

• **Thread cancellation** is the act of terminating a thread before it has completed
  - Example - clicking the stop button on your web browser will stop the thread that is rendering the web page

• The thread to be cancelled is called the **target thread**

• Threads can be cancelled in a couple of ways
  - **Asynchronous cancellation** terminates the target thread immediately
    • Thread may be in the middle of writing data ... not so good
  - **Deferred cancellation** allows the target thread to periodically check if it should be cancelled
    • Allows thread to terminate itself in an orderly fashion
Signal Handling

- **Signals** are used in UNIX systems to notify a process that a particular event has occurred
  - CTRL-C is an example of an *asynchronous signal* that might be sent to a process
    - An asynchronous signal is one that is generated from outside the process that receives it
  - Divide by 0 is an example of a *synchronous signal* that might be sent to a process
    - A synchronous signal is delivered to the same process that caused the signal to occur

- **All signals follow the same basic pattern:**
  - A signal is generated by particular event
  - The signal is delivered to a process
  - The signal is handled by a *signal handler* (all signals are handled exactly once)
Signal Handling

• **Signal handling is straightforward in a single-threaded process**
  - The one (and only) thread in the process receives and handles the signal

• **In a multithreaded program, where should signals be delivered?**
  - Options:
    (1) Deliver the signal to the thread to which the signal applies
    (2) Deliver the signal to every thread in the process
    (3) Deliver the signal only to certain threads in the process
    (4) Assign a specific thread to receive all signals for the process
Signal Handling

• **Option 1 - Deliver the signal to the thread to which the signal applies**
  - Most likely option when handling synchronous signals (e.g. only the thread that attempts to divide by zero needs to know of the error)

• **Option 2 - Deliver the signal to every thread in the process**
  - Likely to be used in the event that the process is being terminated (e.g. a CTRL-C is sent to terminate the process, all threads need to receive this signal and terminate)
Thread Pools

- In applications where threads are repeatedly being created/destroyed thread pools might provide a performance benefit
  - Example: A server that spawns a new thread each time a client connects to the system and discards that thread when the client disconnects

- A thread pool is a group of threads that have been pre-created and are available to do work as needed
  - Threads may be created when the process starts
  - A thread may be kept in a queue until it is needed
  - After a thread finishes, it is placed back into a queue until it is needed again
  - Avoids the extra time needed to spawn new threads when they’re needed
Thread Pools

- **Advantages of thread pools:**
  - Typically faster to service a request with an existing thread than create a new thread (performance benefit)
  - Bounds the number of threads in a process
    - The only threads available are those in the thread pool
    - If the thread pool is empty, then the process must wait for a thread to re-enter the pool before it can assign work to a thread
    - Without a bound on the number of threads in a process, it is possible for a process to create so many threads that all of the system resources are exhausted
Thread-Specific Data

- **Thread-specific data** - in some applications it may be useful for each thread to have its own copy of data
  - May also be referred to as **Thread-local storage** or **Thread-static variables**
  - The `errno` variable is thread-specific

```csharp
// In C#
class FooBar {
    [ThreadStatic] static int foo;
}
```

```java
// In Java
private static ThreadLocal<Integer> threadLocalInt = new ThreadLocal<Integer>();
```

```c
// In a POSIX-compliant system
// see pthread_key_create()
// pthread_setspecific()
// pthread_getspecific()
```
Thread Examples - Windows XP

- Implements threads using the one-to-one thread model

- Also implements a fiber that uses a many-to-many model
  - A fiber is a unit of execution that must be manually scheduled by the application

- Each thread contains
  - A thread id
  - Register set
  - Separate user and kernel stacks
  - Private data storage area
Linux Threads

• **Linux oftentimes uses the term task rather than process or thread**

• **Thread creation is done through the clone() system call**
  - The `clone()` can create either ‘threads’ or ‘processes’ depending on the options passed to `clone()`
    - The options passed to `clone()` determine how much sharing is taking place between the parent and the child
  - A ‘process’ can still be created using the `fork()` system call
  - Provided a specific set of options and the `clone()` and `fork()` systems calls behave identically
Linux Threads (Cont.)

- The following table shows the various flags that can be passed to clone to determine how much sharing is taking place between the parent and the child.

<table>
<thead>
<tr>
<th>flag</th>
<th>meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>CLONE_FS</td>
<td>File-system information is shared.</td>
</tr>
<tr>
<td>CLONE_VM</td>
<td>The same memory space is shared.</td>
</tr>
<tr>
<td>CLONE_SIGHAND</td>
<td>Signal handlers are shared.</td>
</tr>
<tr>
<td>CLONE_FILES</td>
<td>The set of open files is shared.</td>
</tr>
</tbody>
</table>